

PPNPFR12025-03-10

funktionsförflytning  
funktionspredikat: anslutning till funktion  
 $x \mapsto x^2$        $f(x) = x^2$   
 $f := (x \mapsto x^2)$

utvärdering/applikation:  $f(3)$   
 $(x \mapsto x^2)(3)$

verkande i prototyperna.

$\lambda$ -räkning: syntaks  $\frac{\lambda x. e}{e}$   
abstraktformen  $\lambda x. e$   
är en särskaljivit  $x$

odskriften visar att  $f(x) \leftarrow f x$

$\lambda$ -räkningens viktigaste fördel är att den  
är en särskaljivit  $x$ .

$\lambda x. e_1 e_2 e_3 = \lambda x. (e_1 e_2 e_3) \neq (\lambda x. e_1) e_2 e_3$

$$\text{compose} := \lambda g. \lambda f. \lambda x. g(fx) = \\ = \lambda g f x. g f x$$

$$\text{compose } g = \lambda f x. g f x$$

$$\text{compose } g f = \lambda x. g f x$$

$$\text{const} := \lambda c x. c \quad \begin{matrix} \text{use } f \text{ for } t \\ \text{when } c \end{matrix}$$

$$\text{const } + = \lambda x. t$$

Solve nested in program.  
start.

if cero λ izeze true, false in if,  
za kateve velja

$$\cdot \text{ if true } a b = a$$

$$\cdot \text{ if false } a b = b$$

true :=  $\lambda xy \cdot x$

false :=  $\lambda xy \cdot y$

if :=  $\lambda bfe \cdot bfe$

neuen pair: zeigen

- first (pair a b) = a
- second (pair a b) = b

pair :=  $\lambda ab \cdot \lambda f \cdot f \underbrace{ab}_{\text{"object"}}$   
to humans "term reader"

first :=  $\lambda p \cdot p(\lambda ab \cdot a)$

second :=  $\lambda p \cdot p(\lambda ab \cdot b)$

pairer:  
second (pair a b) =  
second (( $\lambda xy \cdot \lambda p \cdot pxy$ ) a b) =  
second ( $\lambda p \cdot p a b$ ) =  
 $(\lambda q \cdot q((\lambda xy \cdot y))(\lambda p \cdot p a b) =$   
 $(\lambda p \cdot p a b)(\lambda xy \cdot y) = (\lambda xy \cdot y) a b = b$

[Chuvchova Fevila]

St. n povedstavimo t izrazom,  
k i spustime fjo f in go u-tcat  
uRovali:

$$0 := \lambda f_x. x$$

$$1 := \lambda f_x. f_x$$

$$2 := \lambda f_x. f(f_x)$$

$$3 := \lambda f_x. f f f_x$$

npr.:  $3 \text{ foo bar} = \text{foo}(\text{foo}(\text{foo bar}))$

naslednje:

n f x.

$$n f x = \underbrace{f(f(\dots f_x))}_{n\text{-tcat}}$$

$$(\text{succ } n) f x = \underbrace{f(f(\dots f_x))}_{n+1\text{-tcat}} = f(n f x)$$

$$\text{succ} := \lambda n f_x. f(n f_x)$$

$$+ := \lambda u w f x. \underbrace{(uf)(wf)}_{\text{steuflr}} x$$

$$\cdot := \lambda u v f x. \underbrace{u(f)}_m x$$

pred/odunit:

$$\begin{aligned} \text{pred} := & \lambda u. \text{second} \left( u \left( \lambda p. \text{pair} \left( \text{succ} \right. \right. \right. \\ & \left. \left. \left. \left( \text{cfirst } p \right) \right) \left( \text{cfirst } p \right) \right) \left( \text{pair } 0 \ 0 \right) \right) \end{aligned}$$

defn si  $f(x, y) = (x+1, x)$ :  
 n-tat applying  $f$  on  $(0, 0)$  is  
 preserving second. + to be pred.

