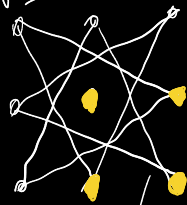


N_3 (graf kupa)



$\gamma G = 4$

$\gamma_t N_3$ ni definirano

→ cikel $C_3 \dots \lceil \frac{3}{3} \rceil = 3$

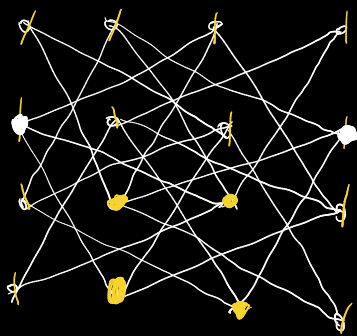
$$\frac{3}{2+1} \leq \delta G \leq \frac{3}{2}$$

$$3 \leq \delta G \leq 4,5$$

spleta G_i

$$\frac{|V_i|}{\Delta G_i} \leq \delta G \leq \frac{|V_i|}{2}$$

N_4 :



$\gamma(N_4) = \lceil \frac{16}{4+1} \rceil = 4$

du: a je možno

$\gamma_t N_4 = 5?$

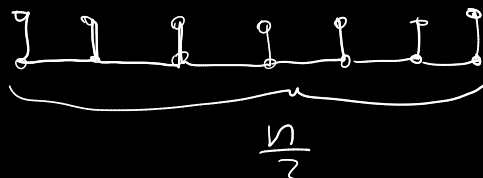
- extension do γ_t
- dominacija in redovna dominacija

N n vozlišč, povezan graf.

a.)

$\delta G = 1$

$\delta G = \frac{n}{2}$



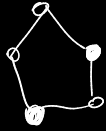
ali pa poti:

$\lceil \frac{n}{3} \rceil = \frac{n}{2}$

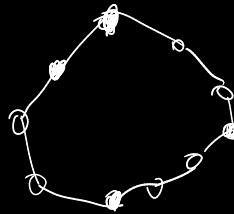
najdi: n vozlišč, povezani

$$\gamma = 2$$

$$\gamma = \frac{2n}{5}$$

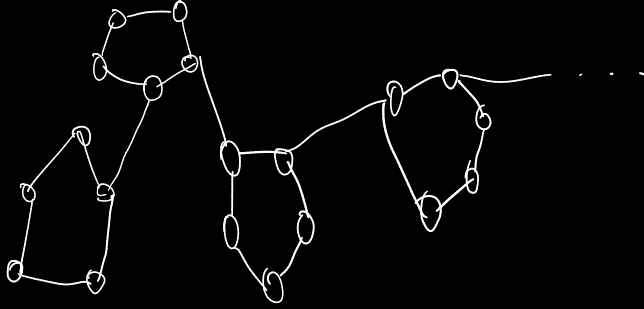


C_5



C_{10}

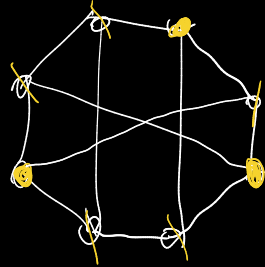
let G_k k -krat C_5 , povezani z mostovi:



N n vozlišč, povezani

c.) $\gamma G = 3$ let $n = 8$ $\gamma = 3$

$$\gamma G = \frac{3n}{8}$$

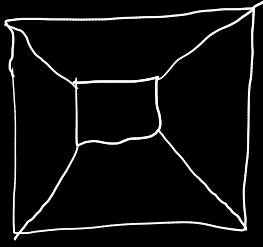


N let G graf, diam $G = 2$ naj daljša razdalja

$$\gamma G \leq \gamma G$$

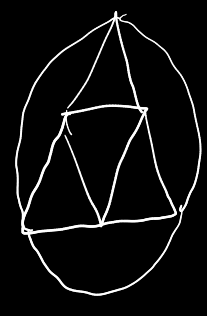
N

Q3:



$$\gamma G \leq \alpha G$$

$$\alpha G \approx G \geq |V G|$$



...

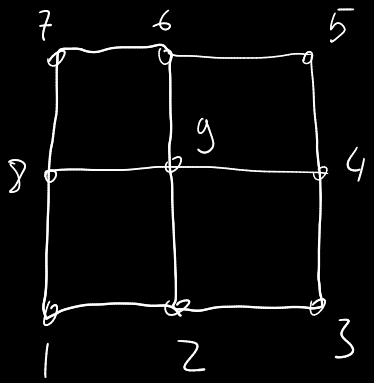
N

zero force

$$P_2 \square P_3$$

imamo 1 operacijo:

- različni in sosedni
sosedom zamenjamo
barvo.



imamo 2 barvi.

našli zaporedje vozlišč, da prebaramo graf iz bele na črno

$$(1, 3, 5, 7, 9)$$

N

✓✓✓

N

Dominacijste igre:

D, S dva igralca
 ↓ ↪ zavladavalka

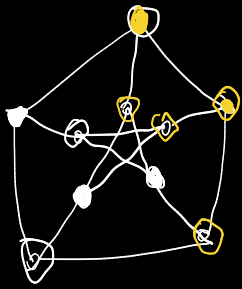
dominator

vsaka poteza: vsaj eno novo vozlišče dominirano
 konec igre, to je izbrana dominantna množica.

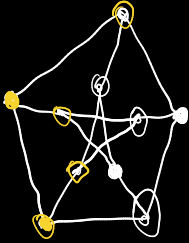
začne D:

γ_g game domination number: najmanjše
 število potez, če oba igralca igrata optimalno

zračne s: γ'_g



$$\gamma_g = 5 = \gamma'_g$$



$$\gamma'_g = 4$$

P_n : $\gamma_g =$



$$\gamma_g P_n = \begin{cases} \lfloor \frac{n}{2} \rfloor - 1; n \cdot 4 = 3 \\ \lfloor \frac{n}{2} \rfloor \text{ sicer} \end{cases}$$

N
 Oskladi, da \exists graf, tjeu je $\gamma_{tg} G < \gamma_g G$?

testo. smo naredili

